

## Problem Set 2

### Problem 1

Let **ALG** be a randomized algorithm with running time  $\mathcal{O}(n^3 + n + \sqrt{n})$  that outputs an optimal solution (for an unspecified optimization problem) with probability at least  $\frac{1}{\sqrt{n \log n}}$ . Give a number  $\ell$  of independent repetitions such that repeating **ALG**  $\ell$  times and returning the best solution results in an algorithm with success probability at least  $1 - \frac{1}{n^7}$ . What is the running time of the resulting algorithm?

### Problem 2

Prove that the **FastCut** algorithm (without repetitions) has a running time of  $\mathcal{O}(n^2 \log n)$ . Is this still true when  $t$  is set slightly larger, for example to  $t := 1 + \lceil (3/4)n \rceil$ ?

### Problem 3

We are given a data stream of numbers  $a_1, a_2, a_3, \dots$  (of unknown length) and want to sample one number  $s$ . However, instead of choosing every item in the stream with the same probability, we want to achieve the following: After seeing  $a_i$ , we want that

$$\Pr(s = a_j) = \begin{cases} 2^{-(i-1)} & \text{for } j = 1 \\ 2^{-(i-j+1)} & \text{for } j \in \{2, \dots, i\} \end{cases}$$

1. To make sure that this defines a discrete probability measure, show that the sum of the desired probabilities  $\sum_{j=1}^i \Pr(a_j)$  after seeing  $a_i$  is always 1.
2. Adapt **ReservoirSampling** such that it stores a number  $s$  which is equal to the different elements in the data stream with the desired probabilities.

### Problem 4

Consider the following recursive and randomized algorithm:

**RandomRecursion**( $\ell$ )

1. Print  $\ell$  on the screen.
2. Toss a random coin.
3. If Heads, call **RandomRecursion**( $\ell + 1$ ).
4. Toss a random coin.
5. If Heads, call **RandomRecursion**( $\ell + 1$ ).

What is the probability that the call **RandomRecursion**(0) finishes running after a finite time (does not run forever)?